Database Theory

Introduction

DBMS - Introduction

- Database Management System (DBMS)
 - "Powerful tool for creating and managing large amounts of data efficiently and allowing it to persist over long periods of time, safely" [UW:01]
- DBMS capabilities
 - Persistent storage
 - Large amount of data; exists independent of any processes
 - Programming interface
 - Easy access and modifications through a *query language*
 - Transaction management
 - Isolation, atomicity, and durability

DBMS - Introduction

Functionalities

- Allows users to create new databases
 - Specification of a schema (logical structure of data)
 - Data-definition language (DDL)
- Allows users to query data
 - Query and modify the data
 - Query language and data-manipulation language (DML), e.g. SQL
- Support the storage of very large amounts of data
 - Secure it from accident or unauthorized use
- Controls access to data from many users at once
 - Prevent to corrupt the data accidentally

DBMS - Introduction

- DBMS vs. File System
 - Conventional *file systems* are inadequate as database systems!
 - Why?
 - They fail to support:
 - Guarantee that data cannot be lost
 - Efficient search and complex queries
 - Efficient modifications to small pieces of data
 - Create a schema (only directory structures)
 - Atomic and independent execution of transactions (concurrent access)

Early DBMS

- Old DBMSs: IMS (hierarchy-based) or CODASYL (graph-based)
- Example: Airline Reservation Systems
 - Data
 - Reservations (flight no, seat, meal preference, ...)
 - Flights (departure, destination, departure/arrival times, etc.)
 - Prices
 - Queries
 - Flights on a given date, time, city, seats available, prices
 - Modifications
 - Booking a flight, assigning a seat, ...
 - Requirements
 - Concurrent access by many agents (make sure that no two agents assign the same seat simultaneously)
 - Avoid loss in case of a sudden system failure

Relational Databases

• Based on the relational model

• Columns: *attributes*

• Rows: *tuples*

- [E.F. Codd. A relational model for large shared data banks. Comm. ACM, 13:6, pp. 377-387, 1971]
- Data is organized as tables, called *relations*

Enrollment	student	course	semester
	Sjöberg	DA2032	Fall 2007
	Jusufi	DA4104	Spring 2008

 Separates the logical view from the physical view (storage structure) of the data

Relational Databases

- Querying a Database
 - Find all students who have taken DA2032 in Fall 2007
 - Structured Query Language SQL

SELECT student FROM Enrollment WHERE course = 'DA2032' AND semester = 'Fall 2007'

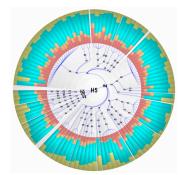
- <u>Answer:</u> Sjöberg, Arlefur, Roos, ...
- So-called *Query Processor* figures out how to answer the query <u>efficiently</u>!

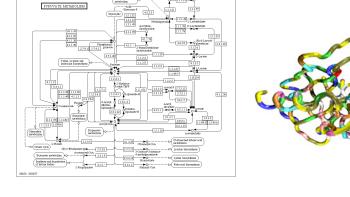
DB in real world

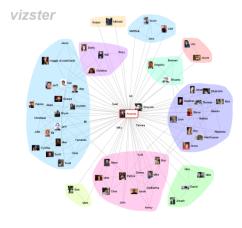
- Relational databases are a great success of theoretical ideas
- Most important DBMS companies are among the largest software companies in the world
 - Oracle
 - IBM (with DB2)
 - Microsoft (with SQL Server, Microsoft Access)
 - ...
- Open-source DBMSs are similarly important
 - MySQL, PostgreSQL, ...
- Rise of NOSQL databases
 - MongoDB, Neo4J, ...

Database Industry and Usage

- Prominent application domains of DBMS
 - Information services on the web, administration, ...
 - Bioinformatics
 - KEGG (biochemical pathways, ...)
 - UniProt (protein sequences, ...)
 - ...
 - Source for Visualization and Visual Analytics
 - Experimental data
 - Statistical data
 - Homeland Security
 - ...







- Smaller and smaller systems
 - In the past, DBMSs could only run on large computers (only large computers could store gigabytes of data)
 - Today, feasible to run a DBMS on a PC
- Bigger and bigger systems
 - Many applications need hundreds of gigabyte
 - Example: A retail chain may store terabytes (10¹²)
 - Newer databases store images, audio, video, etc.
 - Example: Databases for satellite images store petabytes (10¹⁵)
 - Modest size databases stored in *Secondary Storage* (arrays of disks)

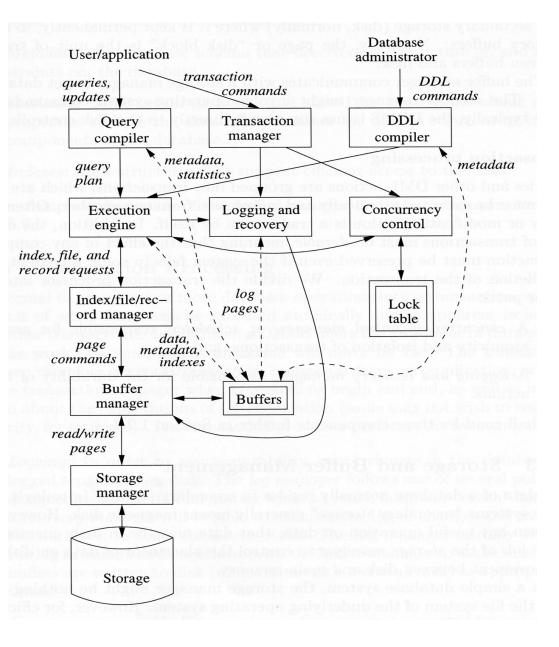
- Dealing with large data
 - What distinguishes database systems from other software?
 - Database systems assume that data is too big to fit in main memory (*Primary Storage*) → data must be located on disk!
 - Today, we need sometimes more than disks!
 - Tertiary Storage
 - Capacity of many terabytes
 - Access takes seconds (ms for disks)
 - CD's, DVD's, etc. may be the storage medium
 - Parallel Computing
 - Satisfies the need of fast access to large data
 - Speeds up through parallel access to disks/devices
 - Needs efficient algorithms to intelligently break queries up

- Client-Server and Multi-Tier Architectures
 - Major database components are on the server, and the client is used to interface with the user
 - In multi-tier architectures, the client is an *application server*, which manages connections to the database, transactions, authorization, etc.
 - Several clients (e.g., web server) \rightarrow application server \rightarrow database server
- Multimedia Data
 - Video, audio, radar signals, satellite images, etc.
 - Challenges
 - Operations (queries, ...) are not simple
 - Users should be able to introduce new functions (object-oriented DBMSs)
 - Query answers can be huge. DBMS must handle such cases

- Information Integration
 - To support information integration, a system should provide a unified view to the database user even if data comes from different sources, for example, different DBMSs of several subsidiary companies
 - Such sources, called *legacy databases*, often use different structures to represent information
 → We need a kind of meta-structure!
 - *Data Warehouses (DWs)* copy information from many legacy databases with an appropriate translation to a central database
 - There are several update strategies
 - DWs are reconstructed each night to reflect updates on legacy databases
- Among other things, this leads to the need of *data mining* techniques, i.e., the search for interesting and/or unusual patterns in data

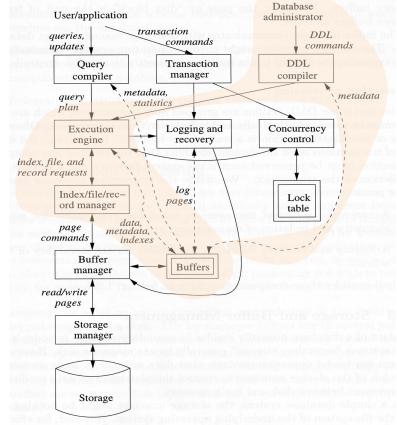
Overview of DBMS

- DBMS Components
 - = system components
 - = data structures
 - = control & data flow
 - = data flow only



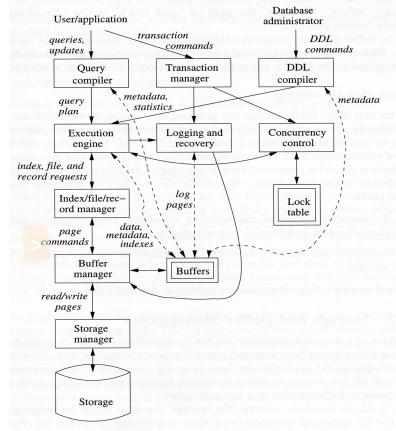
DDL Storage and Buffer

- DDL: Data Definition Language
 - Capable of altering the metadata (schema or structure/constraint information of the database)
 - Requires special authority (i.e., used by database administrator)



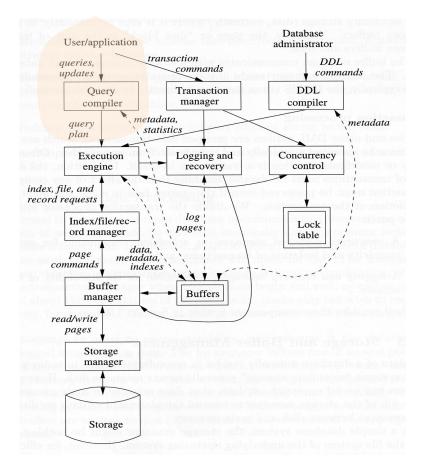
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- Storage and Buffer Management
 - Storage manager controls the location of data on the disk and obtains the blocks containing a file on request from the buffer manager (BM)
 - BM partitions the available main memory into buffers (pages ↔ disk blocks)
 - Information that various components include: data, metadata, indexes (data structures that support efficient access to the data)



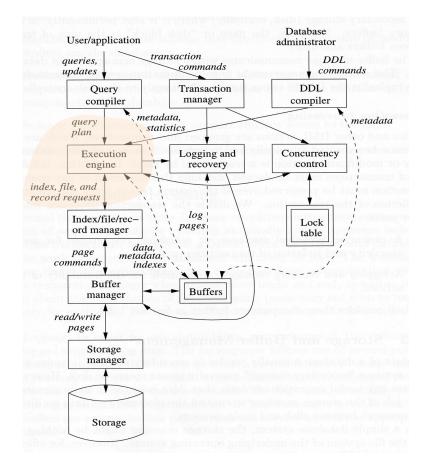
Query Processing

- Query Compiler
 - *Parser* builds a tree structure from the textual form of the query
 - Preprocessor performs semantic checks on the query and translates the parse tree into algebraic operators representing the initial query plan (implementation of a relational algebra)
 - *Optimizer* transforms the initial query plan into the best available sequence of operations on actual data



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- Execution Engine
 - Executes the steps of the chosen query plan
 - Needs to interact with most of the other components of the DBMS



Transaction Processing

- Group queries and/or DML actions into a transaction which must be executed atomically and in isolation from other transactions
- A DBMS guarantees *durability*: the work of a completed transaction will never be lost
- The Transaction Manager performs 3 tasks:
 - Logging
 - Every change in the database is logged on disk (through BM) to enable recovery in case of a crash
 - Concurrency Control assures
 - Atomicity (a transaction is performed either completely or not at all)
 - Isolation (transactions are executed as if there were no other concurrently executing transactions (uses locks))
 - Deadlock Resolution
 - "roll back" or "abort" some transaction

